



Design of Wireless Ad Hoc System using Adaptive Multi-Hop Link State Optimal Routing Protocol

Nayan S. Jambhulkar, Shailesh Kumar, Krushnadeo T. Belerao

Abstract: Now a days for the radio network communication multi-hop routing is used. This multi-hop routing technique covers larger coverage area. Therefore to reach at specific location data is transferred in form of packets from one node to other node. But for the transmission of radio signals over the large distance, large number of transreceivers are required and these transreceivers requires large power to operate. As a result, multi-hop routing can saves energy over separate routing. Therefore it is necessity to design a cost effective multi-hop routing technique for successful transmission of ratio packet data. In this paper a hop by hop adaptive link state optional routing (HALO) is explained. It is the first packet transmitting solution with hop by hop and link state routing, which reduces the cost of transporting data across a packet switch network[3]. The triple model is designed for multi hop packet routing. In this work each node of network iteratively and separately improves the small part of traffic bound. This algorithm finds the shortest path of specific location for every iteration and it is calculated by the marginal cost of the various links of network. The marginal link cost is used to calculate the shortest path between the node and the destination location. This marginal link cost is gathered from link state updates. The various networks changes are automatically identified by the adaptive method which is used in this paper. Due to this the exchange between the packets on wrong node is reduced over the overhead traffic. To validate these theoretical results the experimental evaluations and mathematical calculations are also reported in this work. Net beans java is the programmed use in this proposed project.

Keywords: Link State Advertisement Packets (Lsap), Dijkstra's Shortest Path Algorithm, Link-State Routing, 2 Ack Scheme.

I. INTRODUCTION

Now a day's most of the peoples requires high speed wireless gadgets in different applications. But in wireless network different protocols and standards are utilized and the combination of these protocols are slandered are present in this wireless network. Also in wireless network large number

of heavy software and hardware are used. Therefore it is necessity to design a cost effective wireless system for transmission of data over the network. The transmission of wireless data over the network in form of packets is most suitable and less costly transmission technique used now a days. This packet data transmission technique is generally used for the multipath transmission at the given location and this method consumes less traffic load as compare to other techniques[4]. The optimal routing strategy is used to assign the route so it reduces the cost and traffic of load over the packet switched networks. Over the last 50 years various optimal routing algorithms are developed. But out of these the hop by hop packet transmission by using the protocol distributed link state route is most effective than other technique on the internet.

During the transmission of packet data a specific channel is allocated to it for the transferring and routing the data to the desired location. When the packet data is reached to the desired location then the occupied transmission channel is made ready to other data packet transmission. In this way the transmission channel is reuse for different packet data networks. Payload and header are the two parts of packet. In a payload part software is used which extracts the information present in a packet. The header guides the packet to its location by using hardware.

The hop by hop link state algorithm is the simplest form of algorithm because this algorithm assigns weights at centre to links to control the traffic. This algorithm also submerges the link weights through the network. The link weights calculates the shortest path and forwards the data packet along this path to the desired location. The packet along this path to the desired location. The packets are transferred through the nodes along the network path to the desired location. These are known as intermediate nodes. Switches, routers, and other network hardware devices are widely used as intermediary nodes. fire walls, gate ways, bridges and routers. These routing process and forwarding the data packets are also operated by the computer but it has limited performance.

But in an overall performance of a system, the apparent trade off has been missing. This occurs due to the bad usage of resource over the network with peak traffic. Therefore 30% to 40% network links are only utilizes and others are empty. The main objective of this article is to reduce the trade off among optimality and simplicity of execution in routing. A routing technique Hop by Hop Adaptive link state (HALO) is designed for optimal assignment of routing.

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This is one of the popular and the best routing assignment solution by using link state hop by hop technique. While designing such a solution, there are multiple challenges to conquer. The coordination of various routers only using the link state algorithm is the main challenge to this work. Because of this no router is aware of all other individual communicating pairs inside the network.

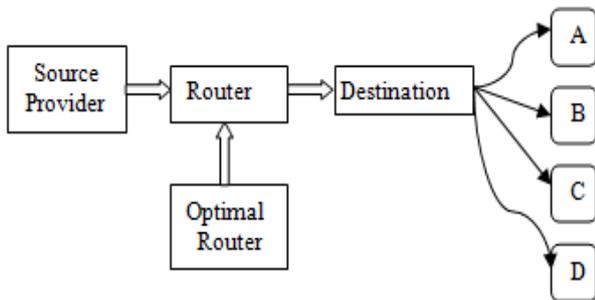


Fig. 1. Hop by Hop Adaptive Link State Optimal Routing (HALO) Routing.

The next challenge is the forwarding of data packets hop by hop technique. In this technique, a router is unable to identify the full path which is emanating by the traffic to its objective. To overcome this problem, the whole path on which the packets take through the network are encoded at the source. For this process a link state algorithm by using gradient technique is used. Figure 1 represents the overview of Hop by Hop Adaptive Link State Optimal Routing (HALO) Routing technique. The network and the input traffic is varies continuously so it is necessity to assign the best route for transferring the packets, this is the another issue emerges. This issue has two sides, in a first case the algorithm requires enough time between the changes of traffic and network to assign and calculate best route. The quasi - static model is fulfill this requirement. In second case the algorithm easily modify the routes. For the sake of clarity of this work, some significant definitions are explained below.

Optimal :- This means minimize total delay. Which is calculated by network operators. Traffic engineering is the task of steering to minimise a particular global cost function, network traffic is routed.

Link - state :- This means every router gets the condition of all network links via flooding link state updates on a regular basis and according to this link state the routing decision made.

Adaptive :- This method in particular adapts and recognize the changes in network such as variations in traffic and changes in topology as infrared signals.

The following is the structure of this paper: The different recommended methodologies employed in this article are presented in section II, which reviews the evolution of the system. The Resources And Techniques Employed are presented in Section III. The simulation of the suggested system as well as the experimental findings are presented in section IV. JAVA Netbeans is used to run the simulation. Finally, part V brings this paper to a close.

II. REVIEW OF THE EXISTING LITERATURE

Asdnibal, Grecian and Romero, investigated the efficient routing issue in autonomous systems, with a focus on optimal

link load balancing for better exploitation of IP network resources. The used theory is part of the MCF (Multicommodity flow) family of network flow issues, and after characterising the sort of alternative given by this type of model, the possibility of implementing it in networks using OSPF and MPLS protocols is examined.

Abdul Hanan Abdullah, Kamalrulnizam Abu Bakar, Khalid Haseeb and Tasneem Darwish, As a result, the primary goal of this study is to introduce the adaptive energy aware cluster-based routing (AECR) protocol, which function to improve energy saving and data service quality. In certain ways, their proposed AECR protocol is different from current energy-efficient routing techniques. To begin, it generated balanced-sized clusters based on node arrangement and avoided the formation of random clusters. It also perfects inter -cluster and intra-cluster routing methods for improving data distribution efficiency while managing data traffic on built-in forwarding routes., and last, it minimises the overdone data traffic on developed forwarding routes. Based on simulation results, the AECR protocol [12] is better than the current state of the art in terms of many presentation measures. The AECR technique divides multi sensor nodes into uniformly sized non-overlapping regions related to network size in this study. As a result, random clusters are prevented, and the duty of CHs is evenly distributed across the overall network region. Furthermore, adopting weighted metrics for the CH electoral system within each cluster region saves time and resources.

The AECR protocol also finds the optimum multi-hop data delivery pathways for the quickest, most energy-efficient, and most reliable data transmissions.

Ozgur Ercetin and Yunus Sarikaya allow access to wireless multihop networks' feasible private information rate areas, where messages are ciphered over large blocks of data. Then, for a particular encoding rate, they devised a dynamic control technique and demonstrated that our algorithm delivers utility arbitrarily close to the maximum attainable value. They then devised a sub-optimal system algorithm for ciphering messages over a certain number of nodes. The simulation results indicate that the methods are efficient and that the suggested algorithm achieves the optimal rates arithmetically. They want to examine distributed versions of their flexible control methods in the future [11], where the planner choice is based on local data. They will also take into account the scenario in which the transmitter only receives poor channel estimation from the connections.

A high-speed but expensive data connectivity network is not always available, particularly in rural areas locations. A store-carry-and-forward-based message carried over current vehicular ad hoc networks is an assured solution to massive data transmission as delay and disturbance tolerant networking under such circumstances (DTN). Masato Tsuru, Agussalim in [10], we explore DTN message delivery situations over many islands where nearby islands are only connected by ferry boat, based on a real-world situation in Indonesia. Messages' source and destination nodes are immobile, off-the-beaten-path islands.

Cars and buses on each island, as well as ferries between islands, relay messages. With some modifications, we extend our previously suggested routing system to multiple-island settings and test its change the characteristics in two delivery direction scenarios.

Masato Tsuru [10] devised an adaptive-spray and hop distance-based technique by refining their previous proposed protocol (A-SnHD).

A-SnHD changes between two phases when a message reaches each island: binary-spray transferring is used for initial dissipation in that island, and subsequently hop distance-based forwarding is utilised in a strict method to prevent superfluous transmission to inappropriate islands.

It has been proved that, despite its simplicity, A-SnHD increases the overall size of delivered messages while lowers the overhead ratio when compared to other methods based on a single simulator-based evaluation in many island circumstances mimicking a real-world situation in Indonesia.

EP is a foundational protocol, whereas PV2 is an artificial protocol that ends with a superior replica limit L in the spray period. Future work should be expected to change depending on local information such as island size. We'll also look into optimising buffer management by adjusting the buffer full condition based on each message's remaining TTL, something they tried in a simpler context in previous work.

Yiftach Richter and Itsik Bergel [9] devised unique wideband WANET routing techniques in which the signal is transmitted in many hops utilising OFDM modulation. They showed how to use a routing system to maximise the routing function. The best route is determined by its local expertise (on the locations of the nodes within its routing zone). They also defined two unsatisfactory low-complexity routing strategies that only use a portion of the knowledge available. These routing strategies come close to matching the performance of the best scheme (especially for low propagation probability, where the quality gap is almost zero). All routing strategies employ decentralised transformation. While they demonstrated that nearest-neighbor routing performs similarly to limited knowledge schemes in some situations (e.g., high transmission probability), they also demonstrated that it performs badly in others (e.g., low transmission probability).

The optimal routing problem is non-convex as a result of these dynamics, and so cannot be solved accurately in general.

Walid Krichene [8] proposed two approaches to approach this non-convex problem: a greedy strategy and a local search method based on the adjoint system, and tested their computational cost and performance using numerical samples.

In particular, he derived adjoint system equations for Hedge dynamics. Their numerical analyses shed light on the tradeoffs and empirical performance of each method: For example, the adjoint approach has the best performance but is quadratic in complexity. The greedy strategy, on the other hand, is easy to implement and performs admirably, despite its limits imposed by its myopic nature. The Hedge adjoint equations' derivation is generic, and it can be used to any optimum control issue involving a selfish population that follows Hedge dynamics.

While its application was limited to the routing game, the

Hedge adjoint equations can be used to solve any optimum control issue involving a selfish population and Hedge dynamics.

Their findings show that the adjoint approach works well in practise, paving the path for further investigation into numerical performance in other applications.

WSNs are distinguishable from conventional ad hoc networks by the limited resources available to sensor nodes. The restricted energy at each node has an impact on network period during data receiving and transmission. In order to achieve the trade-off between energy consumption and data delivery performance, appropriate architecture is required to discover the set of suitable paths for data dissemination while reducing excessive energy consumption among nodes.. Because of their many applications and small installation costs, wireless sensor networks (WSNs) have gained in popularity. The primary goal of a WSN is to decrease energy usage across nodes while ensuring timely and reliable data transmission. Most existing energy conscious routing systems, on the other hand, consume unbalanced energy, resulting in ineffective load balancing and shortened network lifetime.

III. RESOURCES AND TECHNIQUES EMPLOYED

This section depicts the resources and strategies used to put the planned system into action. This section describes how the proposed method evolved.

This approach has been evaluated on a training data set, and several modules are examined in this part when creating a decent routing network model for short-term price forecasting. The block diagram of the suggested system is shown in Figure 2. Whenever the service first begins, the transmitter operates as a source node, sending packets to the destination, which is a receiver, utilising the link state routing principle.

A. Link-State Routing Protocols

Link-state routing protocols are one of the two main types of network topologies used in packet switching networks for computer networking. Distance-vector routing protocols are the other.

The link-state routing protocols System to Intermediate System (IS-IS) and Open Shortest Path First (OSPF) are two examples.

The core idea behind link-state networking is that each node builds a trace of the network's connections in the graph format, indicating which nodes linked to which others.

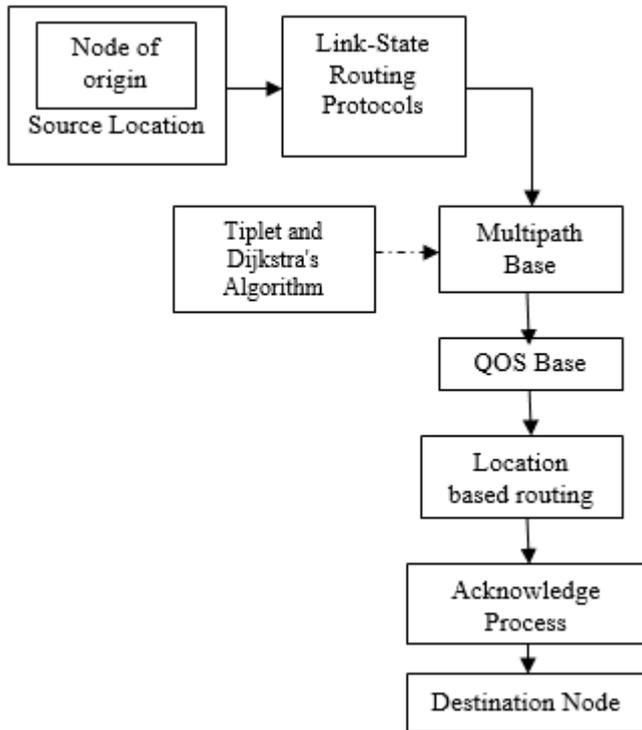


Fig. 2. Block Diagram of Suggested System.

The subsequently best logical path from each node to each probable location in the network is then calculated individually by each node. Each node's routing table will be made up of a collection of the best pathways[5].

The bandwidth is also calculated by the link state protocol based on the network path. By including a minimum BW field in the request packet, the route that has a tendency to break early is discovered and avoided. This minimum BW parameter is used to store a node's available bandwidth. When a node accepts a request packet from a neighbour, it compares the packet's minimum BW value to the available bandwidth on the node[2]. This bandwidth is assigned as the minimum BW if the available bandwidth is less than minimum BW. This procedure will be repeated till the destination is reached. When a destination receives several request packets from various routes, choose the route with the highest minimum BW value and the shortest hop count, and deliver the request packet to the source. That is, we are choosing a route to avoid a node that has a history of increasing delay due to inadequate bandwidth. This Link-State Routing Protocols acquires the packets and then uses The shortest path tree for the destination is computed using Dijkstra's shortest path algorithm.

The following is a step-by-step breakdown of Dijkstra's shortest path algorithm.

Beginning with the source node, Dijkstra's Algorithm analyses the graph to find the shortest path between that node and all other nodes in the network.

The function keeps track of the shortest path between each node and the source node, and it updates these values when a shorter path is found.

The shortest path between two nodes is calculated using the method, and that node is recorded as "visited" and added to the path.

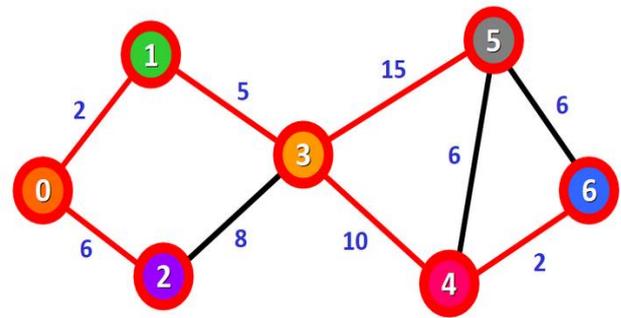


Fig. 3. Example of Dijkstra's shortest path algorithm.

The technique is continued once all of the network's nodes are represented on the path. As a result, we have a path that connects the source node to all other nodes in the most efficient way feasible [3]. The shortest path's edges are indicated by red lines in the diagram. To discover the shortest path to a certain node in the network, you must start at node 0 and follow these edges. Follow the red edges to travel to node 6 from node 0, for example, and the shortest path 0 -> 1 -> 3 -> 4 -> 6 will be automatically followed.

B. Link State advertisement packets (LSAP)

Link State advertisement packets (LSAP) [4] provides the data routers which are needed to develop their databases. Routers do not broadcast their complete routing tables; instead, each router broadcasts only information about routers that are physically close to it. The OSPF routing protocol for the Internet Protocol uses the link-state advertisement (LSA) as a basic communication method. It notifies all other local routers in the same OSPF region of the router's local routing topology. Because OSPF is built for scalability, some LSAs are only flooded out on the interfaces that correspond to the proper region, rather than all of them. This allows for the localization of detailed information while the rest of the network receives summary information.

The nodes in the work are developed using the Java platform. Packets are sent from node one to the last node using the 2ACK protocol, and the last node sends an acknowledgment to the first node. The system then determines whether or not the node is a destination node. If the node is a destination node, it will send an acknowledgment to the sender; otherwise, continue the steps until you reach your destination. For transmitting data from source to destination, the system uses the 2ACK algorithm. The most crucial aspect is that the system sends packets to numerous destinations at the same time, which is referred to as multidestination. We are attempting to reduce the system's energy usage and thereby increase QoS by using Link state routing and the 2ACK method. Using the shortest path technique, the packet is also sent to a multipath destination. On the target path, the 2ACK algorithm forms the triplet module. The nodes that have this triplet use the 2ACK technique to communicate. The 2ACK technique [5] is used to transfer data while lowering congestion by minimising unnecessary acknowledgement. Furthermore, the 2ACK technique ensures that the packets are sent successfully.

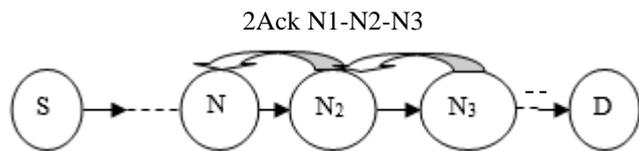


Fig.4: 2 Ack scheme

The 2ACK approach is illustrated in Figure 4. Three successive nodes (triplet) along a path are denoted by N1, N2, and N3. In its Detection Phase, the Link State Routing Protocol builds a direction from a source node S to a destination node D. N1 transmits a data packet to N2, who forwards it to N3, but N1 has no idea whether N3 receives it successfully or not. As a result, the 2ACK technique necessitates an explicit acknowledgment from N3 to inform N1 that the data packet has been successfully received. When node N3 successfully receives a data packet, it sends a 2ACK packet to node N1 as well as the data packet's ID across two hops (i.e., in the opposite direction of the routing path as shown in fig.). The 2ACK packet recipient is marked by N1 in the triplet [N1N2N3], while the 2ACK packet sender is designated by N3 [3].

C. Configuration of nodes in the network

Destination node packets are pumped into the system on a regular basis to build the connections between the various networks. Request messages must be responded to by routing reply messages to the source. Each node that does not respond to messages is removed from the network. The source form takes a circuitous path to the destination, passing through intermediary nodes.

D. Routing with several paths

Multipath routing, which uses many pairs of paths among a source and a destination to fulfil these criteria, is a promising routing technique. Multipath routing is a network routing approach that uses numerous alternative paths to provide a assortment of benefits such as greater bandwidth and security. The numerous pathways computed may overlap, be edge disjointed, or have nodes that are disjointed from one another.



Fig.5: The suggested Adaptive multi-hop link state optimum routing in WANet is designed as a system.

Multipath routing techniques have been the subject of extensive investigation[7].

IV. RESEARCH OUTCOMES

This chapter entails the performance of the developed wireless ad hoc network using Adaptive multi-hop link state optimal routing protocol. First, the training data must be collected. The training begins through one input nodes and gradually increases as the network's performance improves.

The method for the present scheme is shown below.

- 1) Get started.
- 2) The sender is S, and the recipient is D.
- 3) Using the HALO approach, compute the shortest path tree for destination D.
- 4) Using three nodes from the shortest travel route, create a new triplet N1 N2 N3.
- 5) N1 is a temporary sender node, N3 is a temporary receiver node, and N2 is a temporary middle node.
- 6) From N1 to N3, the packet is transmitted.
- 7) Using 2Ack, N3 sends an acknowledgement to N1.
- 8) Determine whether (N3==D).
- 9) If not, repeat Steps 4–8 for the next three nodes.
- 10) If the answer is yes, D sends S a final acknowledgement.
- 11) Stop.

Figure 5 shows the suggested multi hop adaptive link state optimum routing in WANet is designed as a system. The GUI for send file, Received History, Send History, Received File, Graphs, Show graphically Route and compression are design in this system as shown in figure 5.

Figure 6 shows the output of the Adaptive multi-hop link state optimum routing protocol system used in the proposed wireless ad hoc network.

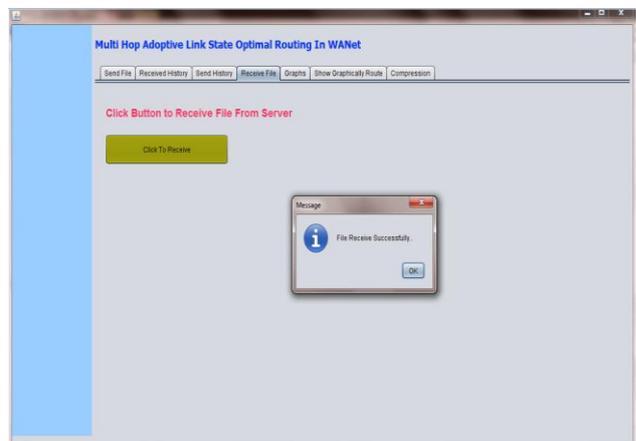


Fig.6: The output of the Adaptive multi-hop link state optimum routing protocol system used in the proposed wireless ad hoc network.

It is evident from this diagram that the file sent by the sender was successfully received. Figure 7 shows the performance parameters of the Adaptive multi-hop link state optimum routing protocol system used in the proposed wireless ad hoc network. The different parameters such as File Name, File Size, Received Data, Received Time, File Received From and Suggested Path are calculated.

Fid	File Name	File Size	Received Date	Received Time	File Received From	Shortest Path To Re.
0	Distance Vect...	10.220 MB	11/06/2017	06:00:28 PM	mttch-PC	
1	saadip.jpg	30MB	9/1/2017	12:30AM	PC3	
2	sdids.jpg	10MB	9/1/2017	10:30AM	PC3	

Fig.7: The Performance Parameters of The Adaptive Multi-Hop Link State Optimum Routing Protocol System Used in The Projected Wireless Ad Hoc System

V. CONCLUSION AND IMPLICATIONS FOR THE FUTURE

The proposed system comes to a conclusion based on the outcomes of the proposed model. The experimental findings suggest that the received files through the network are fairly accurate. The JAVA Netbeans programme was utilised to produce this proposed work. The purpose of this paper is to reduce the amount of the tradeoff in routing between optimality and employment comfort. Using the shortest path algorithm, the packet must be sent to the multipath destination. It provides an optimal approach and a method for sending packets to several destinations (multidestination). The packets will arrive at their location in a timely manner. Furthermore, traffic engineering is efficiently performed, resulting in traffic control and thus congestion control methods. This paper also includes error-controlling strategies that are effective. The suggested strategy is intended to serve many destinations while maintaining the maximum level of QoS in terms of energy consumption. These properties make data transit from source to destination more efficient.

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