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Abstract: Any organization is obliged to ensure secrecy of data from hacking criminals complying with the increasing demand for secured data. So data preservation is indispensablethrough cryptographic methods. It is adoptedseveralreal life applications such as ecommerce, In this paper we indicated a procedure for generating different Pythagorean triples and with the help of C++ coding developed a mechanism for both encoding and decoding of a plain text in English alphabets. We also demonstrated them with illustrative examples.

Keywords: Pythagorean Triple, Encoding, Decoding, Wireless Sensor Networks (WSN).

I. INTRODUCTION

Is is a challenge to devise an efficient procedure for key generation. The author in [MonishaPrabhu. (2013)] proposed one such scheme through Primitive Pythagorean triples (PPTs) as pointed outin [Kak (2010)]. In [ArtanLuma and BujarRaufi. (2014)] the authors suggested a Pythagorean Triple Algorithm. Using this they extended the definition of the Pythagorean Theorem. That is, there exist two or sometimes even three solutions to Pythagorean Theorem for any two numbers p and q such that either $p \in 2Z$ and $q \in 2Z+1$ or vice versa. Based on these they created an encryption and decryption key for simple symmetric cryptosystem.

It is interesting to know that there are many methods available to obtain Pythagorean triples (PT). In [Duvvuri Surya Rahul and SnehanshuSaha (2015)] the authors highlight a structured approach to find all PPTs with no repeat digit in the triples. A connection between the sum of side lengths and prime numbers are observed of a primitive right triangle. A Fundamental result about Primitive Pythagorean Primes (PPP) on the lines of Fundamental Theorem of Arithmetic is conjectured. One cannot find a variety of methods in the literature to identify all PPTs and PTs with no repeat digit. In [Barning (1963)] the author developed a nice procedure for determining such triples. Price [Price (2008)] later came out with his own method.One can also refer to [Saunders and Randall (1994)], Mitchell's formula [Mitchell (2001)]and [William (2004)] for more on this

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There are several analytical and numerical methods [Mushtaque and Snehanshu (2014)] for estimating prime numbers. Cryptographic algorithms are divided into two types viz., symmetric-key and public key algorithms. In symmetric-key algorithm, for both encryption and decryption same key is adopted and it has be sent via a communication channel or sometimes manually. A drawback is that a data thief may get an accessto the secret key through illegal means during transmission. Symmetric key algorithms are further divided into stream cipher and block cipher algorithms. Symmetric-key encryption algorithm through block cipheris less preferred than to a cipherstreamon smaller units of plaintext, usually bits or bytes. Stream ciphers are much faster than block ciphers. Most of the stream ciphers facing the problem of generating one random bit in each round of process as the output stream of cryptosystem [Majid Bakhtiari and MohdAizainiMaarof. (2011)]. This increases of algebraic correlation against cryptosystems[Majid Bakhtiari and MohdAizainiMaarof. (2011), Meier and Staffelbach, (1990)]. Vernam's one-time pad is one of the stream cipher involving a string of randomly generated bits. As the whole keystream is random, an opponent can guess the plaintext. It is noted that though, Vernam's onetime pad is perfectly secured, remembering and storing such a key is too tedious because the size of the key is always taken as the size of plaintext and hence it is at least practical [Charles Pfleeger and Shari Lawrence Pfleeger. (2007)]. In DES symmetric-key algorithm, initially 64-bit key called DES-key is given as input and in each round a 56-bit round key has been generated to produce the ciphertext. Suppose, if the DES keys chosen are weak keys, the Hamming distance between a plaintext and a ciphertext produced by DES may be less giving the data thief a chance to recover from the ciphertextthe intended plaintext. To overcome these, a PPT based key stream can be generated and its length can be determined by both the sender and the receiver as in the case of Vernam's one time pad. Further to develop from the PPT based keystream a new DES key, the crucial seed PPT and beginning position should be known by the participating parties and this aspect makes the keystream from the key sequence unpredictable. This makes the the PPT based keystream better than traditional DES as in the former different key is generated for each round and in the latterin all rounds only one key is used.

II. PYTHAGOREAN TRIPLES

Generation of events of with given probability is pertinent in cryptography [Kolmogorov (1965), Kak, and Chatterjee (1981), Kak (1985), Kak (1987) Kak (1989), Kak (2007)].



One can produce infinite number of Pythagorean triples of the form x (a, b, c) where x>1 and such triples have the ability of creating probability events. A primitive Pythagorean triple (a, b, c) is the one in which a, b and c are pairwise relatively prime. For a detailed description see [Kak (2010),Euclid (1997), Heath (1956)].

Kak (1987), Kak (2011), O'Conner and Robertson, (2000), Seidenberg (1978), Parakh and Kak(2009). Parakh and Kak(2011)].

III. PYTHAGOREAN TRIPLES FOR ENCRYPTION AND DECRYPTION

An ordered triple $(a,b,c) \in \mathbb{Z}^3$ is called a Pythagorean triple if $a^2+b^2=c^2$. One familiar approach to generate a Pythagorean triple is: for any x, $y \in Z^+$ with x>y,set $a=x^2-y^2$, b=2xy, $c=x^2+y^2$. This formula is due to Euclid. Then Newton opined an integer solution: a is d times (x^2-y^2) , bis 2d times xy,c is d times (x^2+y^2) with x, y>0, x > y,gcd(x,y)=1, x and y differ in parity and d is the gcd of (a,b,c). We call the triple (a,b, c) to be Primitive Pythagorean if d=1. Suppose that $a^2+b^2=c^2$ and gcd(a,b)=1. Then there exists $r,s \in Z$ such that c=a+r, c=b+s, where gcd(a, r)=1 an gcd(b,s)=1. Clearly a+r=b+s and a-s=b-r. Set a-s=b-r= μ . So that a=s+ μ and b=r+ μ . Then we get $c=r+s+\mu$. So $(a=s+\mu, b=r+\mu, c=r+s+\mu)$ standsfor new solution to the system $a^2+b^2=c^2$. That is, $(r+\mu)^2+(s+\mu)^2=(r+s+\mu)^2$. This gives μ^2 =2rs.Here one can pick r and s in such away that we deduce a solution of the type $s=2x^2,r=y^2,s>r,gcd(x,y)$ =1. Putting these values in μ^2 =2rs, we get μ^2 =4x²y² and $a=2x^2\pm 2xy$ we hence $\mu=\pm 2xy$.Hence get $b=y^2\pm 2xy$, $c=2x^2+y^2\pm 2xy$. This reveals a fact that for any two given integers x and y, with either $x \in 2Z$ and $y \in 2Z+1$ or vice versathere are (a_1,b_1,c_1) and (a_2,b_2,c_2) that acts as solutions.

Example 1: If x=3,y=1 then we have two solutions $a_1=24,b_1=7,c_1=25$ and $a_2=12,b_2=-5,c_2=13$.

We now use the above said procedure to encrypt and decrypt a text. If 't' stands for plaintext, α for the key and e for the encrypted message then e=t+ α (mod 26) where the alphabets a to z are represented by 0 to 25, in order.To decrypt the above we have to perform t=e- α (mod 26) and the key is generated by using the values.x₁=2x²+2xy,x₂=2x²-2xy, x₃=2xy;y₁=y²+2xy, y₂=y²-2xy, y₃=x²-y²; z₁=2x²+y²+2xy, z₂=2x²+y²-2xy, z₃=x²+y² and (x₁,y₁, z₁), (x₂, y₂,z₂), (x₃,y₃,z₃) is reduced with mod 26 and create an encryption key of the type: x₁, y₁, z₁, x₂, y₂, z₂, x₃, y₃, z₃.

Example 2: Suppose that we have SASTRA DEEMED TO BE UNIVERSITY as plaintext.

Table.1 shows the message encoding for the above word.

Table 1 Message Encoding

Tuble I litebouge Encouring															
S	A	S	T	R	A	D	E	E	M	E	D	T	0	В	E
1	0	1	1	1	0	3	4	4	1	4	3	1	1	1	4
8		8	9	7					2			9	4		

U	N	I	V	E	R	S	I	T	Y
20	13	8	21	4	17	18	8	19	24

We encrypt using x=7,y=2. Simple calculation reveals x_1 =2.7²+2.7.2=126; y_1 =2²+2.7.2=32; z_1 =2.7²+2²+2.7.2=130; x_2 =2.7²-2.7.2=70; y_2 =2²-2.7.2=-24;

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Table 2 Message Encryption															
S	A	S	T	R	A	D	E	E	M	E	D	T	0	В	E
1	0	1	1	1	0	3	4	4	1	4	3	1	1	1	4
8		8	9	7					2			9	4		
2	6	0	1	2	2	2	1	1	2	6	0	1	2	2	2
2			8		2		9		2			8		2	
1	6	1	1	1	2	5	2	5	8	1	3	1	1	2	6
4		8	1	9	2		3			0		1	6	3	
0	G	S	L	T	W	F	X	F	Ι	K	D	L	Q	X	G
U	Ν	I	V	E	R	S	Ι	T	Y						
2	1	8	2	4	1	1	8	1	2						
0	3		1		7	8		9	4						
1	1	2	6	0	1	2	2	2	1						
9		2			8		2		9						
1	1	4	1	4	9	2	4	2	1						
3	4					0		1	7						
N	0	E	В	E	J	U	E	V	R						

Note that Newton's formula for generating Pythagorean triples sometimes duplicates Pythagorean triples.For instance, if d=1,x=9,y=3;d=9,x=3,y=1; andd=18, x=2, y=1 results in (a,b,c)=(72,54,90);(a,b,c)=(54,72,90). In literature only a few methods exists that produce distinct Pythagorean triples.

Table 3 Message Decryption \mathbf{o} G \mathbf{S} L Т W \mathbf{F} X \mathbf{F} Ι D \mathbf{L} Q K X G 2 5 2 3 1 6 1 1 1 5 8 1 1 1 2 6 9 2 4 8 1 3 0 6 3 2 0 1 2 2 2 1 1 2 0 1 2 2 2 6 6 2 8 2 8 1 0 0 3 4 4 1 1 1 1 1 1 1 8 8 9 7 9 2 \mathbf{S} \mathbf{S} T A R A \mathbf{E} \mathbf{E} D \mathbf{T} D \mathbf{E} M В N 0 \mathbf{E} В \mathbf{E} J U \mathbf{E} V R 4 4 2 2 1 1 1 9 4 1 3 4 1 7 1 1 2 0 1 2 2 2 1 6 9 2 8 2 2 1 8 2 4 1 1 8 1 2

9 4

Algorithm for Encoding

Ι

Algorithm ENCODE(List)

This algorithm reads the input and prints a report. Pre:List contains the plain text entered by the user. Post:Cipher text will be displayed.

7 8

E R S

1.if(List not equal to ".")

1.read List

2.increment count

2.end if

0 3

3.decrementcount, set count1 to count.

4.loop(count1 not equal to 0)





1.convert list elements to respective ascii values and set to Algorithm for Decoding convertedlist. Algorithm DECODE(List) 2.decrement count. This algorithm reads the input and prints a report. 5.end loop 6.set randomvalue1 to value1. Pre:List contains the cipher text. 7.set randomvalue2 to value2. Post:Plain text will be displayed. 8.setlistP to set1 pythagarean values. 1.if(List not equal to ".") 9.setlistQ to set2 pythagarean values. 1.read List 10.setlistR to set3 pythagarean values. 2.increment count 2.end if since, onlyine values..... 11.loop(index<3) 3.decrementcount, set count1 to count. 1.if(element of listP>=0) 4.loop(count1 not equal to 0) 1.setlistP element to mod listP element. 1.convert list elements to respective ascii values and set to 2.else convertedlist. 1.incrementlistP element by 26. 2.decrement count. 3.end if 5.end loop 4.if(element of listQ>=0) 6.set randomvalue1 to value1. 1.setlistQ element to mod listQ element. 7.set randomvalue2 to value2. 5.else 8.setlistP to set1 pythagarean values. 1.incrementlistQ element by 26. 9.setlistQ to set2 pythagarean values. 10.setlistR to set3 pythagarean values. 7.if(element of listR>=0)since, only nine values..... 1.setlistR element to mod listR element. 11.loop(index<3) 1.if(element of listP>=0) 8.else 1.incrementlistR element by 26. 1.setlistP element to mod listP element. 9.end if 2.else 10.increment index. 1.incrementlistP element by 26. 12.end loop 3.end if since, only 3 sets 9 with elements..... 4.if(element of listQ>=0) 13.setlistC to all the sets elements respectively. 1.setlistQ element to mod listQ element. 14.set count2 to count. 15.set index to 0 1.incrementlistQ element by 26. 16.loop1(count2 not equal to 0) 6.end if 1.if(index < 9)7.if(element of listR > = 0) 1.newlist[count]=convertlist[count]+listC[index] 1.setlistR element to mod listR element. 8.else 2.else 1.incrementlistR element by 26. 1.set index to index mod 9 9.end if 2.newlist[count]=convertlist[count]+listC[index] 10.increment index 12.end loop 4.increment index since, only 3 sets 9 with elements..... 5.increment count 13.setlistC to all the sets elements respectively. 6.decrement count2 14.set count2 to count. 17.end loop1 15.set index to 0 18.loop3(count not equal to 0) 16.loop1(count2 not equal to 0) 1.if(newlist element >25) 1.if(index < 9) 1.decrementnewlist element with 26 1.if(convertedlist[count]>=listC[index]) 2.end if 1.newlist[count]=convertlist[count]-listC[index] 3.decrement count 19.end loop3 1.newlist=convertlist[count]-listC[index]+26 20.set choice to userchoice 3.end if 21.if(choice is equal to UPPERCASE) 1.incrementnewlist to 65 1.set index to index mod 9 2.convertnewlist elements to respective ascii characters. 2.if(convertedlist[count]>=listC[index]) 3.displaynewlist 1.newlist[count]=convertlist[count]-listC[index] 22.else if(choice is equal to LOWERCASE) 1.incrementnewlist to 97 1.newlist=convertlist[count]-listC[index]+26 2.convertnewlist elements to respective ascii characters. 4.end if 3.displaynewlist 3.end if 4.increment index 23.else 1.go to statement 20 5.increment count 24.end if 6.decrement count2 end ENCODE 17.end loop1 18.set choice to userchoice



- 19.if(choice is equal to UPPERCASE)
 - 1.incrementnewlist to 65
 - 2.convertnewlist elements to respective ascii characters.
 - 3.displaynewlist
- 20.else if(choice is equal to LOWERCASE)
 - 1.incrementnewlist to 97
 - 2.convertnewlist elements to respective ascii characters.
 - 3.displaynewlist
- 21.else
- 1.go to statement 18
- 22.end if
- end DECODE

C++ Program For Encoding

```
int p[3],q[3],r[3],x,y,co=1,n,i,j,flag,d[100],k=0;
char a[1000],b[100];
cout<<*enter the letters with a '.' at the end:'<<endl;
for(i=0;a[i=1]!='.';i++)
(cin>>a[i];
0;a[i]='.';i++)
    or(i=0;i<n;i++)
if((int)a[i]>96)
         b[i]=((int)a[i]-65);}

**onter_v=2
                nter value1,value2(x,y)";
cin>>y;
or(i=0;i<3;i++)
if(i==0)
i] = ((2*x*x)-(2*x*y));
i] = ((y*y)-(2*x*y));
i] = ((2*x*x)+(y*y)-(2*x*y));
  for(i=0;i<3;i++)
{if(p[i]>=0)
```

```
se
=q[i]+26;
(r[i]>=0)
i]=r[i]%26;
      [10];
p[0];c[1]=q[0];c[2]=r[0];c[3]=p[1];c[4]=q[1];
r[1];c[6]=p[2];c[7]=q[2];c[8]=r[2];
=0;4<9;i++)
<c[i]<<end1;
   ;
(i=0;i<n;i++)
 (j<9)
i]=b[i]+c[j];
  =(j%9);
i]=b[i]+c[j];
  r(i=0;i<n;i-
f(d[i]>25)
i]=d[i]-26;
 i]=d[i];
it flag1=0;}
          enter ur choice as either l(lower) and u(upper) case
       you have been entered a wrong option change it. ';
       h;
endl;]
     <<endl;
i=0;i<n;i++)
      cout<<setw(2)<<(char)a[i]<<' ';
      cout<<endl;
=0;i<n;i++)
out<<setw(2)<<(int)b[i]<<' ';
cout<<endl;
for(i=0;i<n;i++)
```

```
;
cout<<endl;
```

```
A 8 T R A D E E M E D T O B E U N I V E R S

0 18 19 17 0 3 4 4 12 4 3 19 14 1 4 20 13 8 21 4 17 18

0 18 11 10 22 5 23 5 8 10 3 11 16 23 6 13 14 4 1 4 9 20

17

0 8 L T M P X P I K D L Q X O N O E B E J U
```

IV. DISCUSSION ON PYTHAGOREAN TRIPLES AND ITS ROLE IN WSN

The C++ program codes given below makes use of Pythagorean triples, where the input words given in text form by the user are converted into corresponding numerical values. According to Newton's formula for each set of two different numbers there exists three sets of Pythagorean triples. For Encoding: The input information which are converted into numerical value are manipulated according to the formula $t = e-\alpha \pmod{26}$ using Pythagorean triples. Then numerical values are displayed as alphabets either as upper case or as lower case as per the user choice. For Decoding: We first consider the numeric value that corresponds to the entered code. Then by adding the manipulated set of Pythagorean triples to the numerical values as per the condition where necessary, we adjust the numerical values to the alphabetical number range to decode the input information.

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One can see in the C++ code that alphabets entered are converted into corresponding numerical values. Then user will exercise his choice of the set of two numbers to generate three sets of Pythagorean triples. Then by manipulating the numerical values through the Pythagorean triples one can derive the encoded input. Similarly, to decode the encoded information one can convert the entered alphabets to corresponding numerical values and with the help of generated Pythagorean triples read the message.

One can view in literature how to produce efficientkeys for instance [Blom (1985), or Chan et.al (2003), Eschenauer and Gligor (2002)] who take into account variety of factors such as energy, memory and cost etc., The generation of Pythagorean triples also finds immense potential for application in developing key management schemes in WSN. WSN are vital for smart environments to achieve automation in facilities such as transportation, home, utilities and industries. For this we look to sensory data from real world. We depend on distributed WSN for collecting such data and we process it to elicit crucial information. With the threat of stealing, we intend to minimize loss on the nearby nodes and the network. To ensure this messages exchanged between two nodes should adhere to proper encryption and authentication. To begin a communication a sender and the receiver must have a prior knowledge about common key. Pythagorean triples are involved in good depth to derive certain novel approaches such as symmetric-key algorithms for generating the key from a key stream. The main advantage of this is that the key value generated from the key stream is chosen by both the sender and the receiver. Also the key sequence size is arbitrary in length. As the generated key stream is random both the sender and the receiver need not remember such keys. We hope to revert more on this elsewhere.

V. CONCLUSION

We have given a crisp introduction about the Pythagorean triplet generation mechanisms available in the literature and then indicated a new procedure to obtain different sets of Pythagorean triples for the same set of input values. Then by using the Number theory formulae we have outlined a procedure to encode and decode any word given in English alphabets.

C++ Program for Decoding



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