

# Compact Lightweight Cryptographic Algorithm for Optimization of Resources

Sunitha Tappari, K. Sridevi, Durga Rao Jenjeti

**Abstract:** The study of Lightweight cryptography has been one of the interesting topics in symmetric cryptography in the recent years. Lightweight symmetric ciphers has gained interest due to the increasing demand for security services in constrained computing environments, such as in the Internet of Things (IoT). Though the protocols of light weight in are providing more security in various applications, resource utility is more in key generation, key scheduling, permutation layer and substitution box layer operations. More resource usage make possible to have high utilization of power and overhead of area. In this paper, a novel method is proposed to decrease the utility of resources, which follows register reutilization scheme. The resource reutilization is coordinating with the delay in Substitution Permutation Network (SPN). Our design is compared, using area, throughput, power, and latency as metrics.

**Keywords:** Symmetric light weight cryptography, Resource optimization, Register reutilization, Cryptography coding.

## I. INTRODUCTION

The world over in the environment computing can be embedded. Nowadays, the broadband Internet is generally accessible to any user and its cost of connectivity is also reduced, more sensors and gadgets are getting connected to it [1]. There are many research works focusing on complexities around the IoT [2]. For each valuable transmitting data, the sophisticated sensors are implanted in the substantial equipments adjacent to us. The devices itself start the sharing process of large data which should communicate with IoT with more security.

The IoT is taking the mobile network, conventional internet and sensor network to a different level are connected to the internet. IoT is further susceptible to serious threats of security than available network, the reason for that is it includes resource constrained objects, assorted nature, open environment employment, lively behavior. Security in IoT environment should ensure accurate implementation of confidentiality, integrity, authentication, non-repudiation, and access control [3].

A key element of secure communication of the system is Cryptography [4]. Conventional cryptography is not suitable for resource constraint devices. The majority of these devices use limited power sources to the point where it is required to rely on energy harvesting [5], power optimization techniques [6], and novel transmission technologies [7]. Therefore it is difficult to provide cryptographic solutions for constrained environments. Lightweight Cryptography (LWCRYPT) is a new region of research that will reach the requirements place by smart devices. In this, cryptography methods have to work with minimum amount of vital resources of required objects. "Lightweight" term represents algorithm with less consumption of energy, and less computational power requirement. LWCRYPT is a cryptography technique, customized for less resource requirement devices.

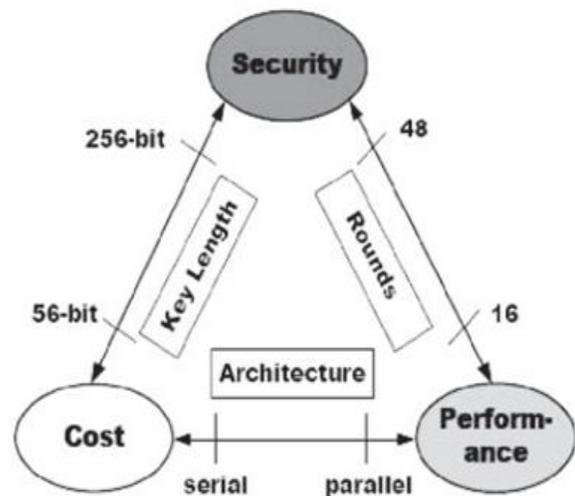


Fig.1: Scaling measures among Performance, Security and Cost. The key length is decreased to 56-bits from 256-bits; the processing rounds count is reduced to 16 from 48 in the logic of encryption, the mode of architecture shifts from parallel to serialized as indicated in Fig.1. In addition, requirement memory is decreased to Kilo bytes from Giga bytes and processing speed comes down from GHz to KHz.

One of the classification of LWCRYPT technique [8] is symmetric cryptographic algorithms Vs asymmetric cryptographic algorithms. Due to the hardware-friendly nature, Symmetric cryptography is suitable for constrained devices. Which implies minimization of supplementary computational cost and consumption of power.

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II. PRESENT CIPHER

PRESENT is a symmetric light weight block cipher. In hardware realizations of PRESENT [9], its circuit size is small and that enables implementation in the RFID tag, which is not possible using the standard AES encryption. PRESENT is hardware oriented block cipher algorithm of lightweight cryptography. It can serve with the input of 64bit blocks and available keys of 80/128 bits [10]. The cipher is dependent on a Substitution-Permutation Network (SPN), with a round-based processing system. As a result of its straightforwardness, it is commonly utilized in circumstances where low-control utilization and high chip proficiency are wanted. This design is derived directly from the algorithm specification and a top level description of PRESENT is shown in Fig.2.

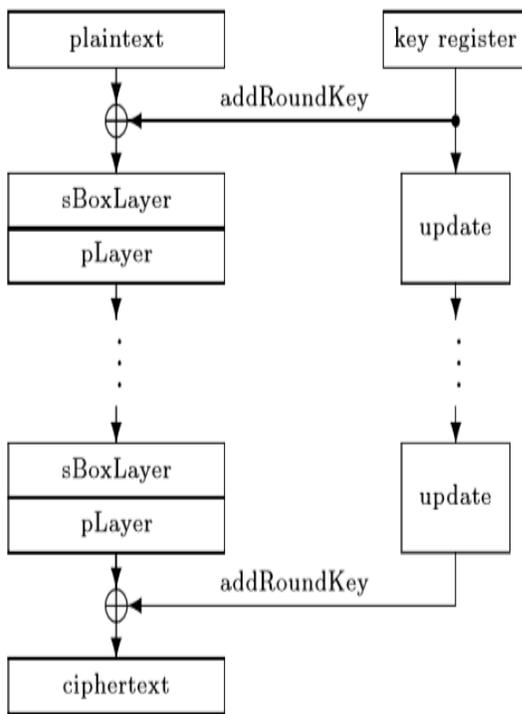


Fig.2. A top-level algorithmic description of PRESENT.

The resource overhead is generated in this approach due to recursive operation is more in key scheduling approach. This approach is proficient, but it requires large quantity of resources. To diminish the resource overhead, approach of resource minimization is proposed in which delay of resource is tuned.

III. IMPROVED PRESENT CIPHER

The improved PRESENT cipher overcomes the resource overhead. The customized structure of the resource controlled light weight architecture is shown in Fig.3.

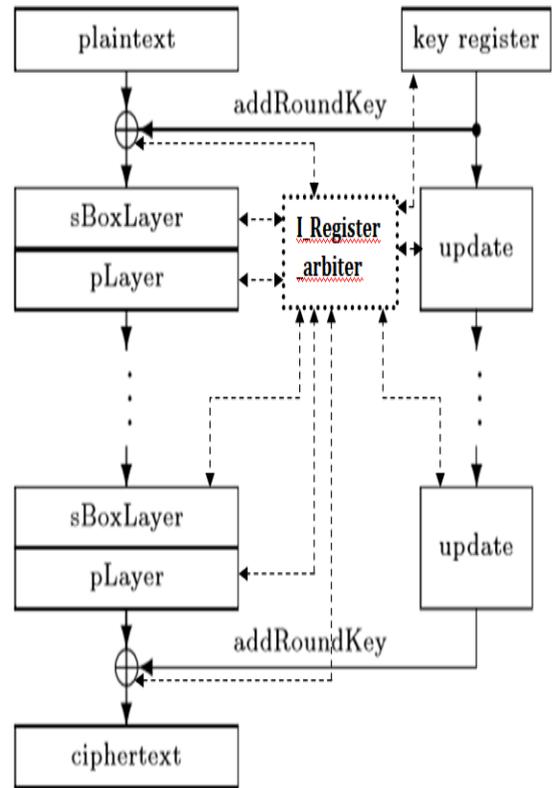


Fig.3. A top-level resource controlled light weight PRESENT.

In this approach, a 4 operational instructions set is used. The flow of data in those instruction operations are shown in Fig.4.

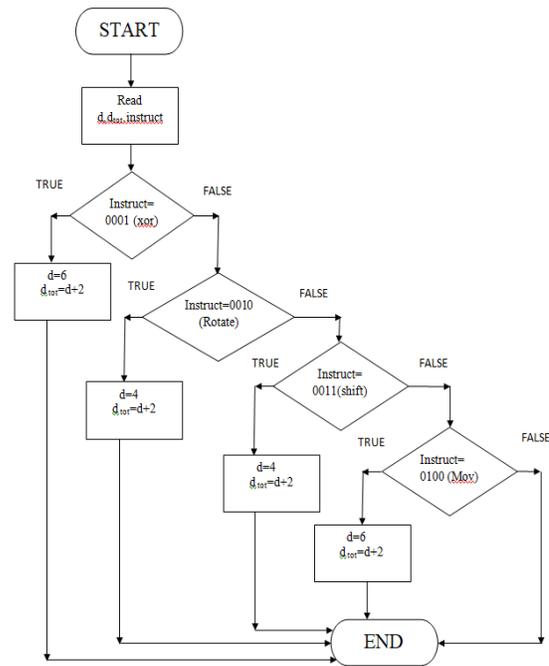


Fig.4 Data flow for instruction set.

To execute each XOR instruction, it requires, 6 delay cycles which includes 1 decode clock cycle, 2 computational clock cycles and 3 register allocation clocks. For ROTATE instruction, it requires, 4 delay cycles which includes 1 decode clock cycle, 1 computational clock cycles and 2 register allocation clocks. For SHIFT instruction, it requires, 4 delay cycles which includes 1 decode clock cycle, 1 computational clock cycles and 2 register allocation clocks. For MOV instruction, it requires, 6 delay cycles which includes 1 decode clock cycle, 2 computational clock cycles and 3 register allocation clocks. Also, each instruction to be fetch and write for operation. Resource employment in regular instruction category will reduce the overhead.

In this approach, the common clocking instructions are realignment as a set of instruction and a common delay in register allocation is defined. Single time POST coding operation is performing by this realignment process, and the delay is previously mentioned as a refery array. In this case, the executed operations do not go off for calculation of clock delay and the delay buffer reference. Fig.5 represents Data flow for updated instruction set.

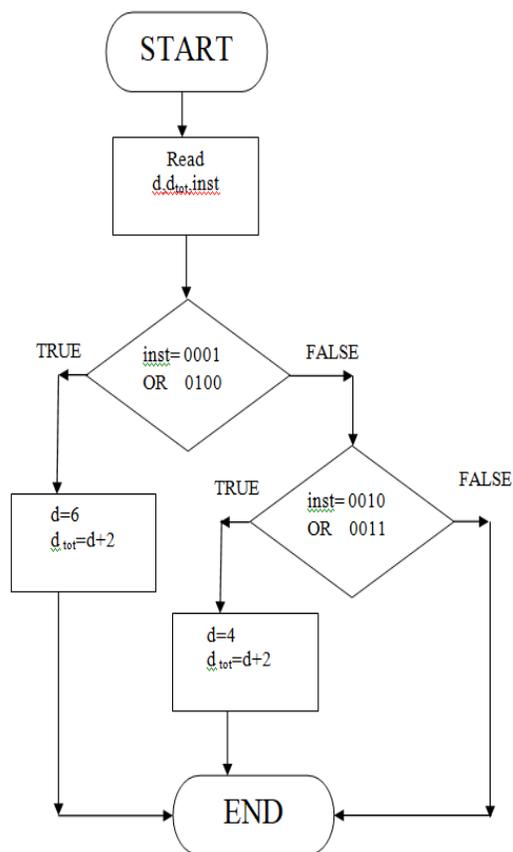


Fig.5 Data flow for Updated instruction set.

This approach reduces the latency of clock, thus the huge requirement of resource computations are eliminated. The proposed clock allocation process has an allocation of delay values for each instruction precomputed. To give best resource

utility, these instructions are programmed in cryptography coding.

#### IV. SIMULATION RESULTS

A Hardware description language for the proposed approach is developed and synthesized on Xilinx ISE tool. The timing observations are simulated in same tool. FPGA design is chosen for the realization of the improved approach.

The latency, power, area and throughput metrics are evaluated. For the test appraisal, 4 instruction set of operations are performed to carry out key updating, substitution and permutation processes. The alignment of registers for a set of instruction is resulted in timing waveforms is shown in Fig.6.

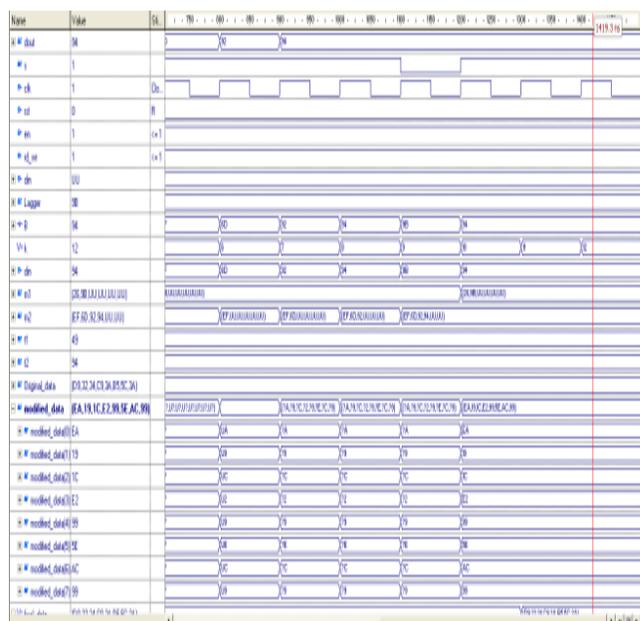


Fig.6 Alignment of registers for a set of instructions

The optimized approach in security coding is observed and as outlined below,

In designed approach, a 64-bit block of data and 80-bit key is considered as

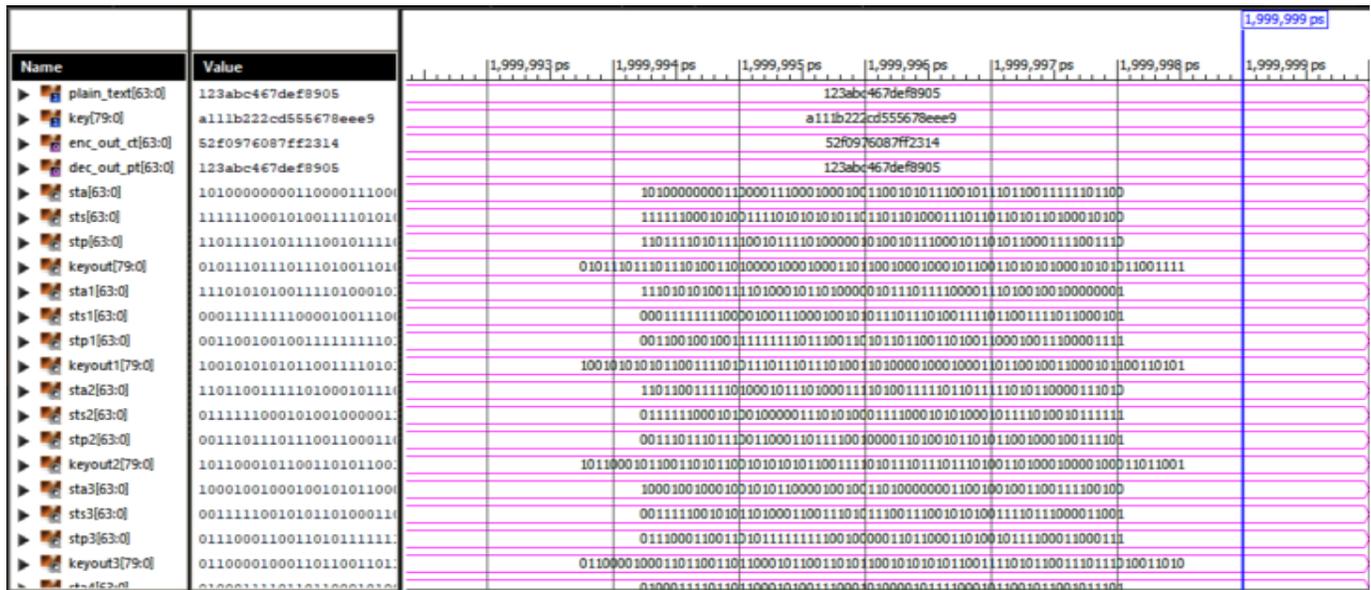
- Plain Text (PT):  
"123ABC467DEF8905"
- Key(K):  
"A111B222CD555678EEE9"
- Encrypted output(ENC\_out):  
"52F0976087FF2314"

In encryption process, plain text is encrypt with the key, it will generate cipher text (Encrypted output). In decryption process, the cipher text is acts as a input and it is decrypt with the same key, it will generate plaintext text (original data).

The Obtained Decrypted output data is given as

- DEC\_out\_PT:  
"123ABC467DEF8905"

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**Fig.7. The Waveforms for Encryption process and Decryption process**

From Fig.7, the Data is recovered same as Input data (Plaintext). The implementations of FPGA is targeting to Xilinx FPGA device Spartan-3 (xc3s200-5ft256). The synthesis report For the targeted device, is illustrated in Table.1

**Table 1: Synthesis Implementation for targeted FPGA.**

Device Utilization Summary			
Logic Utilization	Used	Available	Utilization(%)
• Number of Slices	99	1,920	5%
• Number of Slice Flip Flops	156	3,840	4%
• Number of slice LUTs	174	3,840	4%
• Number of occupied Slices	114	1,920	5%
• Number of Slices containing only related logic	114	114	100%
• Total Number of 4 input LUTs	174	3,840	4%
• Number used as a route thru	128		
• Number of bonded IOBs	149	173	86%
• IOB flip flops	14		
• Number of GCLKs	1	8	12%
• Total equivalent gate count for design	1,643		
• Additional JTAG gate count for IOBs	721		

To estimate, implementation, performance, Spartan-3 is targeted are valid as follows.

The power is measured by using Xilinx Xpower analyzer tool. The system latency is tested for a targeted FPGA device and it is diminished by resource reutilization method. The max. throughput is a function of system frequency ( $F_{max}$ ), a block latency cycles and the block size is computed as,

$$THR = \frac{F_{max} \times B_{size}}{LAT} \quad (1)$$

The area is represented in terms of no. of flip flops, no. of LUTs and no. of slices.

The matrix comparison of existing work [11] and the improved work is illustrated in Table.2.

**Table:2 The comparison of different parameters.**

Parameter	Existing work (PRESENT)	Improved work (IPRESENT)
• Latency (cycles)	133	129
• Power (mW)	271.5	192
• Throughput (Mbps)	102.89	188.22
• Area (FF+LUT+SC)	492	428

## V. CONCLUSION

This paper presented the improvement of cryptography coding with optimal resource utilization. The new technique enhances the operation speed by using limited playing field calculation method to diminish the calculation time for Encryption and decryption of data using PRESENT. The surveillance from the realization of improved PRESENT approach, the device can support to 64-bit data for modified with a rate of 379.4 MHz. The throughput is about 188.22 Mbps. The resource occupied on Spartan-3 targeted chip is about 543 required slices. The optimized resource utility encryption is achieved and observed.

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